

### **AOSTA**

## and some ideas

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## Overview

Part I:

AOStA Revisited: A Short Overview

Part II:

A Code Generator for AOStA

Part III:

**Beyond Performance** 





#### Part I: AOStA

- AOStA: Adaptive Optimizations for Smalltalk

Profiled execution to identify hot spots

Compiles to optimized bytecode

Dynamic deoptimization (debugging)

- Written in Smalltalk





### More...

#### Profiled execution: two areas for JIT-compiled methods

- The optimized area works as usual
- In the unoptimized area methods have a counter for each send and backward branch

#### **Collecting type information**

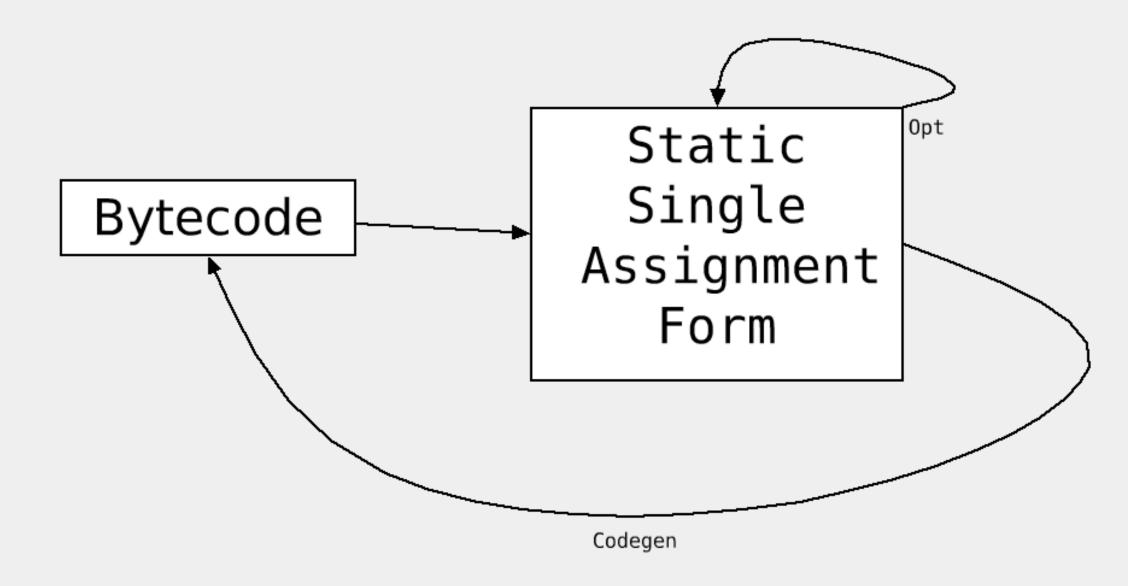
- via Polymorphic Inline Caches
- need to be readable from Smalltalk

#### **Optimizations:**

e.g., Inlining
Specializations for known types



# Bytecode-to-Bytecode







### **Status**

#### 2003:

- Design
- Frontend: Bytecode transformed to SSA
- Middle: SSA Framework, sample optimizers
- For VisualWorks

#### 2004:

- Backend: transformation out of SSA
- Simple Code Generator
- Done with Squeak





### Part II: More about the backend

- Short introduction to SSA (Static Single Assigment)

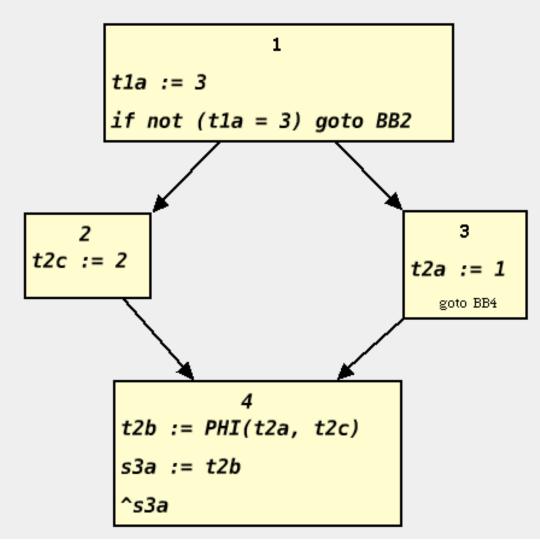
- Two steps:
  - 1) Deconstruction of SSA
  - 2) Code generation

- Some examples



## SSA - Static Single Assigment Form

- SSA: Each Variable has **one** assignment
- If controll flow merges, we need to select the variable from the path we came from







### SSA

- Very nice for many optimizations
- but: Code generation not possible directly
  - --> Need to remove virtual selector functions (PHI-functions)

Two step code generation

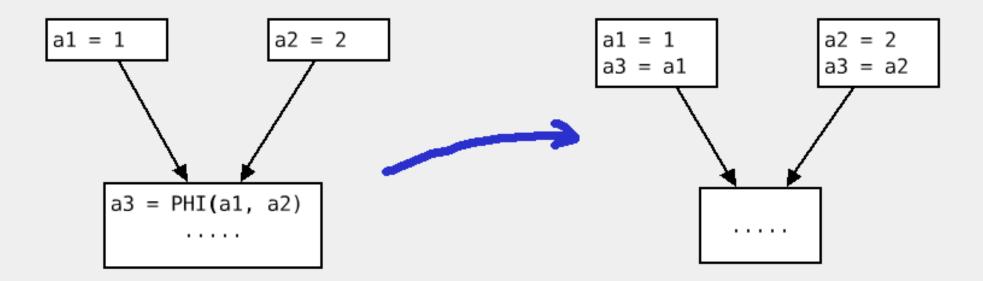
- 1) Deconstruction of SSA
- 2) Code generation





### SSA Deconstruction

#### Canonical method:



#### Problems:

- Wrong after some optimizations
- Copies need to be removed





## Phi-Congruency Method

Method by Vugranam C. Sreedhar, Roy Dz-Ching Ju, David M. Gillies, und Vatsa Santhanam.

**Idea:** Transform program that all variables are the same in PHI:

$$a1 = PHI(a1,a1)$$



$$a1 = a1$$

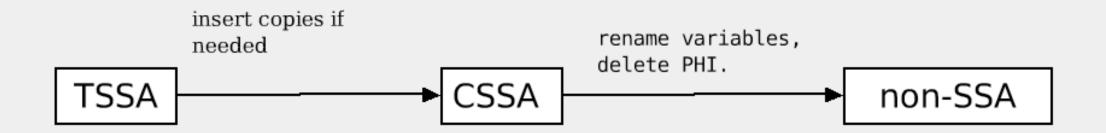
- Insert copies
- Renaming





## Phi-Congruency: Overview

- Two step process:



- Nice properties:
  - Without any optimizations, no copies are needed
  - Simple heuristics for copy placement

### Number of copies

canonical:  $\sim 16000$ 

new: ~1000 (without Opt:





### **IRBuilder**

### Symbolic Assembler

### **Example:**

#### Execute:

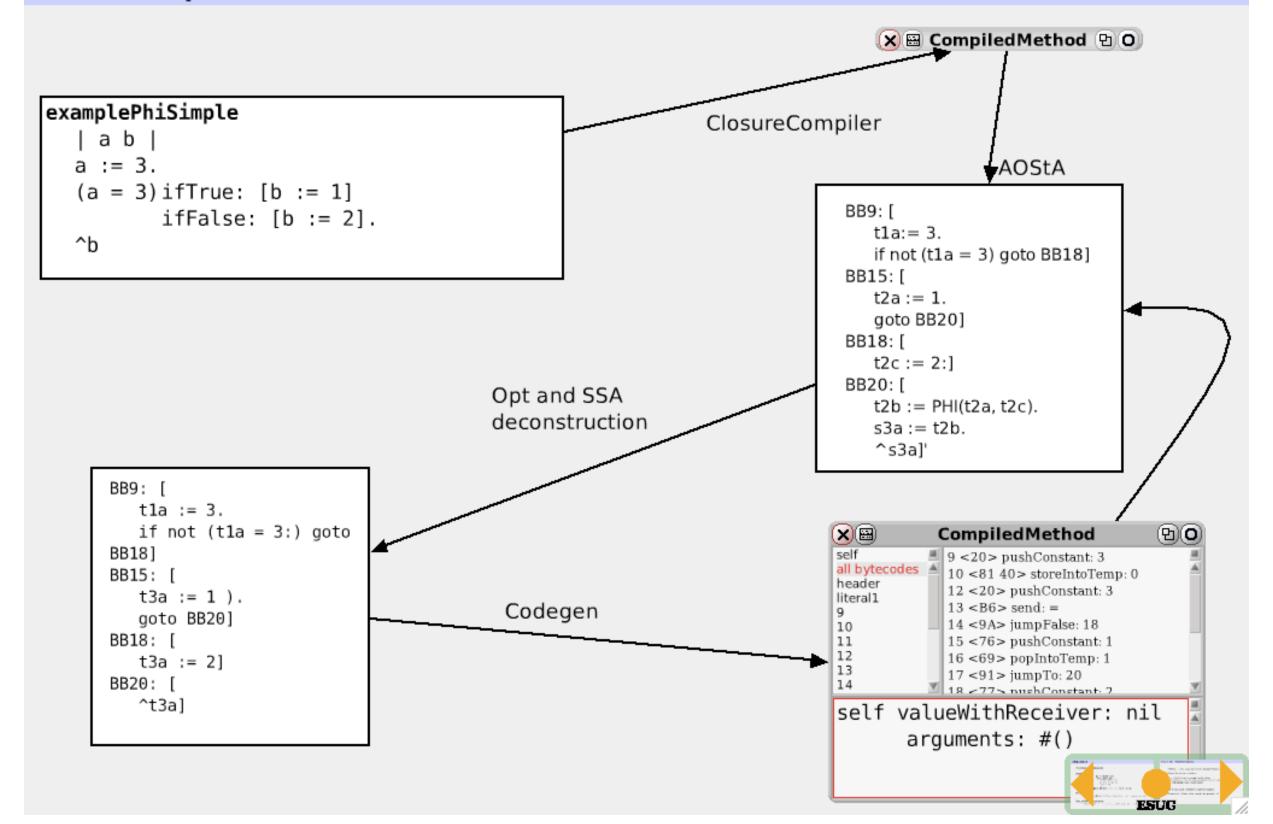
```
aCompiledMethod valueWithReceiver: nil arguments: #()
```

### Install in the system:

```
Float addSelector: #test withMethod: aCompiledMethod.
```



## Example





## Part III: TODO/Ideas

TODO... lots. e.g. dynamic deoptimization

Possible experiments:

- -> AOStA on Squeak with Jitter
- -> Does it make sense with just an interpreter?
- -> Exupery as a backend

All these are related to performance.

Question: What else could be possible?



## Runtime Translation as a System Service

- Enables more late binding

### Example:

- iVars are accessed via offsets
- offsets are calculated at compile time
- makes changes and experiments harder

Make a MOP practical

- Allows a much simpler System





## MOPs and other strange stuff

MOP: Meta Object Protocol.

Idea: Provide an API for changing the language semantics and implementation at runtime.

(e.g., meaning of inheritance)

For Squeak: MetaClassTalk

- Nice, but slow
- A runtime translator could regain performance

Example: ClassBoxes



## Two Kinds of Bytecode

"Image"level Vs. Interpreter Level

- Imagelevel bytecode can be simple:
  - => No optimizations at all

- Imagelevel bytecode and interpreter bytecode could even be different:
  - => Latebinding of the execution format
- Why not just use the AST?



## "2 Worlds"

### Software-Engineering

- AST instead of Bytecode
- late bound
- no optimizations

# Translator

#### Execution

- bytecode or binary
- optimized
- late binding resolved

